

# EIGENVALUES AND OPTIMIZATION

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# 1. OUTLINE

## PART I: SYMMETRIC MATRICES

- Hyperbolic polynomials, Lax, and convex optimization
- Convex spectral functions and unitarily invariant norms
- Duality, subgradients, and some Lie algebra

## PART II: NONSYMMETRIC MATRICES

- Spectral radius, transient dynamics, and Kreiss's theorem
- Pseudospectral analysis and optimization
- Distance to uncontrollability

**PART I:**

**SYMMETRIC MATRICES**

## 2. HYPERBOLIC POLYNOMIALS

**Example:** Consider the homogeneous polynomial

$$p(u, v, w) = u^3 - 2uv^2 - uw^2 + 2v^2w.$$

For all real  $v, w$ ,  $p(u, v, w) = 0 \Rightarrow u$  **real**.

We say  $p$  is **hyperbolic relative to**  $d = (1, 0, 0)$ :  
 $t \mapsto p(x - td)$  always has all real roots.

**Why?**

$$p(u, v, w) = \det \begin{bmatrix} u & v & w \\ v & u & v \\ w & v & u \end{bmatrix}.$$

**Lax Conjecture (1958)** Hyperbolic polynomials on  $\mathbf{R}^3$  relative to  $(1, 0, 0)$  are all of the form

$$p(u, v, w) = \det(uI + vA + wB) \quad \text{with } A, B \text{ symmetric.}$$

**True:** Helton/Vinnikov 2002, Lewis/Parrilo/Ramana 2004.

### 3. HYPERBOLICITY CONES

Hyperbolic polynomials

- are simply defined;
- are common (there are open sets of such polynomials);
- have surprising convexity properties.

**Theorem (Gårding 1951)** Hyperbolic  $p$  relative to  $d$  has convex **hyperbolicity cone**—component of  $d$  in  $\{x : p(x) > 0\}$ .

Linear optimization over hyperbolicity cones is **tractable!**

**Example:** The **determinant** is hyperbolic relative to  $I$  on

$$\mathbf{S}^n = \{n \times n \text{ real symmetric matrices}\} :$$

each  $X \in \mathbf{S}^n$  has all real eigenvalues  $\lambda_1(X) \geq \dots \geq \lambda_n(X)$ .

The hyperbolicity cone is  $\mathbf{S}_{++}^n = \{\text{positive definites}\}$ .

Hence **semidefinite programming** (generalizing LP).

## 4. CONVEXITY AND SYMMETRY

Convexity of  $\mathbf{S}_{++}^n$  and  $-\log \det$  are special cases of:

**Theorem (Davis 1957)** Convexity and permutation-invariance of  $f : \mathbf{R}^n \rightarrow \overline{\mathbf{R}}$   $\Rightarrow$  convexity of

$$X \in \mathbf{S}^n \mapsto f(\lambda_1(X), \dots, \lambda_n(X)).$$

(Eg:  $f(x) =$

$$\begin{cases} 0 & (x > 0) \\ +\infty & (x \not> 0) \end{cases} \quad \text{or} \quad \begin{cases} -\sum_i \log x_i & (x > 0) \\ +\infty & (x \not> 0). \end{cases} )$$

Extends to hyperbolic  $p$  (relative to  $d$ ): interpret  $\{\lambda_i(x)\}$  as the roots of  $t \mapsto p(x - td)$  (Bauschke... 2001). Similarly:

**Lidskii's Theorem (extended)**  $\lambda(z) - \lambda(x)$  is a convex combination of permutations of  $\lambda(z - x)$ .

Lax conjecture reduces both results to  $\mathbf{S}^n$  case (Gurvits 2004).

## 5. INVARIANCE

A function  $F : \mathbf{S}^n \rightarrow \overline{\mathbf{R}}$  is **spectral** if

$$F(U^T X U) = F(X) \quad \text{whenever} \quad U^T U = I$$

because then the spectral decomposition  $\Rightarrow$

$$F(X) = F(\text{Diag}(\lambda_i(X))).$$

The Davis result **characterizes** convex spectral functions:

**Theorem** A spectral function  $F$  is convex  $\Leftrightarrow F$  is convex on  $\mathbf{D}^n = \{n \times n \text{ real diagonals}\}$ .

Reminiscent of...

**Theorem (von Neumann 1937)** Unitarily invariant  $G : \mathbf{M}^n \rightarrow \mathbf{R}$  is a norm  $\Leftrightarrow$  the restriction  $G|_{\mathbf{D}^n}$  is a norm.

## 6. DUALITY

Von Neumann's proof depended on a **duality formula**:

$$G_*|_{\mathbf{D}^n} = (G|_{\mathbf{D}^n})_*,$$

for the **dual function**

$$G_*(Y) = \sup\{\langle X, Y \rangle : G(X) \leq 1\}.$$

We can prove Davis's result by analogy. The **Fenchel conjugate** of a function  $F : \mathbf{E} \rightarrow (-\infty, +\infty]$  is

$$F^*(Y) = \sup\{\langle X, Y \rangle - f(X)\},$$

and the duality formula

$$F^*|_{\mathbf{D}^n} = (F|_{\mathbf{D}^n})^*.$$

holds for spectral  $F$ .

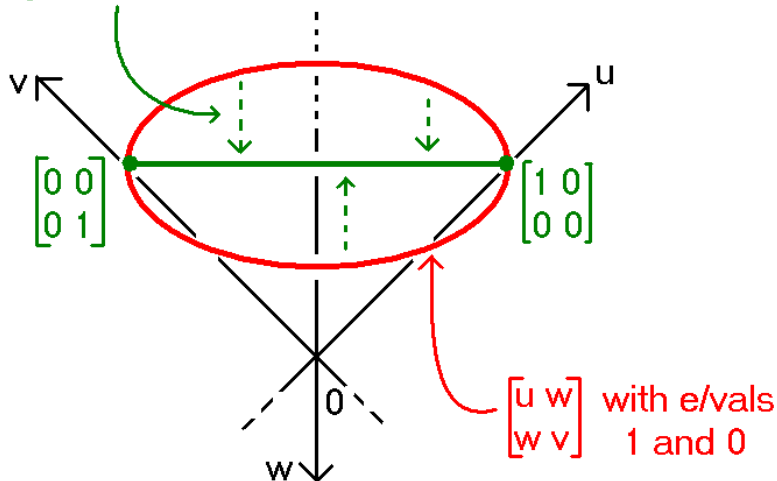
What's the unifying thread?

## 7. HORN'S THEOREM (1954)

One route to the Davis result...

**Theorem** Convex combinations of permutations of  $y \in \mathbf{R}^n$  give all possible diagonals of  $X \in \mathbf{S}^n$  with eigenvalues  $\{y_i\}$ .

projecting onto  
{diagonals}



## 8. AN ALGEBRAIC FRAMEWORK

**Theorem (Kostant 1973)** Consider

- a real semisimple Lie group  $G$ , with Lie algebra  $\mathfrak{g}$ ;
- a maximal compact subgroup  $K \subset G$ , with Lie algebra  $\mathfrak{k}$ ;
- the corresponding **Cartan decomposition**  $\mathfrak{g} = \mathfrak{k} \oplus \mathfrak{p}$ ;
- a maximal abelian subspace  $\mathfrak{a} \subset \mathfrak{p}$ .

Then for  $x \in \mathfrak{a}$ ,  $\text{proj}_{\mathfrak{a}}(K \cdot x) = \text{conv}(W \cdot x)$ ,  
where  $W$  is the **Weyl group**.

**Example** For Horn's theorem, take

$$\begin{aligned} \{U \in \mathbf{M}^n : U^T U = I, \det U = 1\} &\subset \{U : \det U = 1\}, \\ \{\text{skews}\} \oplus \{\text{traceless symmetric}\} &= \{\text{traceless}\} \\ \mathfrak{a} = \{\text{diagonals}\} &\text{ and } W = \{\text{permutations}\}. \end{aligned}$$

## 9. CONVEXITY AND INVARIANCE

The von Neumann and Davis results extend to this setting:

**Theorem** Given a maximal compact subgroup  $K$ , and a maximal abelian subspace  $\mathfrak{a}$  of a corresponding Cartan subspace  $\mathfrak{p}$ , consider an **invariant** function  $F : \mathfrak{p} \rightarrow \overline{\mathbf{R}}$ :

$$F(k \cdot x) = F(x) \quad (k \in K, x \in \mathfrak{p}).$$

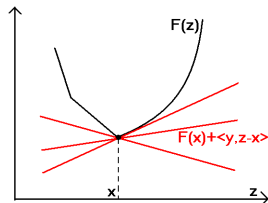
Then we have:

**Convexity characterization**  $F$  is convex  $\Leftrightarrow F|_{\mathfrak{a}}$  is convex.

**Duality formula**  $F^*|_{\mathfrak{a}} = (F|_{\mathfrak{a}})^*$ .

**Convex subdifferential**  $y \in \partial F(x)$ ,

i.e.  $F(x) + \langle y, z - x \rangle \leq F(z) \quad \forall z \in \mathbf{E}$



$$\Leftrightarrow y = k \cdot v, \quad x = k \cdot u, \quad \text{with } k \in K, \quad v \in \partial F|_{\mathfrak{a}}(u).$$

In the  $\mathbf{S}^n$  case, this becomes...

## 10. SPECTRAL SUBDIFFERENTIALS

**Theorem** If  $f : \mathbf{R}^n \rightarrow \overline{\mathbf{R}}$  is convex, permutation-invariant, so  $X \in \mathbf{S}^n \mapsto F(X) = f(\lambda(X))$  is convex (by Davis), then:  
 $Y \in \partial F(X) \Leftrightarrow$  **simultaneous spectral decomposition**,

$$U^T U = I, \quad U^T (\text{Diag } x) U = X, \quad U^T (\text{Diag } y) U = Y,$$

for some  $U$  and  $y \in \partial f(x)$ .

The subdifferential extends to nonconvex, Lipschitz  $f$ :

$$\partial f(x) = \text{conv} \{ \lim \nabla f(x_r) : x_r \rightarrow x \}$$

(Clarke, 1974). The same characterization holds.

An illustration of nonsmooth analysis in linear algebra. . .

## 11. EIGENVALUE PERTURBATION THEORY

**Theorem (Lidskii 1950)** If  $X, Z \in \mathbf{S}^n$ , then  $\lambda(Z) - \lambda(X)$  is a convex combination of permutations of  $\lambda(Z - X)$ .

### A proof via nonsmooth analysis

- Via a separating hyperplane, we need, for any  $w \in \mathbf{R}^n$

$$w^T(\lambda(Z) - \lambda(X)) \leq [w]^T \lambda(Z - X),$$

where  $w \mapsto [w]$  maps components into decreasing order.

- Consider the (nonconvex) spectral function

$$F(X) = w^T \lambda(X).$$

A **nonsmooth mean value theorem** shows

$$F(Z) - F(X) = \langle Y, Z - X \rangle$$

for some  $Y \in \partial F(W)$  where  $W \in [X, Z]$ .

- Now apply the subdifferential formula. □

**PART II:**

**NONSYMMETRIC MATRICES**

## 12. THE SPECTRAL RADIUS

**Question** How should we design a (parametrized) square matrix  $A$  to force  $A^n \rightarrow 0$  quickly as  $n \rightarrow \infty$ ?

**Theorem** Rate of decay

$$\inf\{\mu : A^n = O(\mu^n) \text{ as } n \rightarrow \infty\}$$

equals **spectral radius**

$$\rho(A) = \max\{|\lambda| : \lambda \in \Lambda(A)\},$$

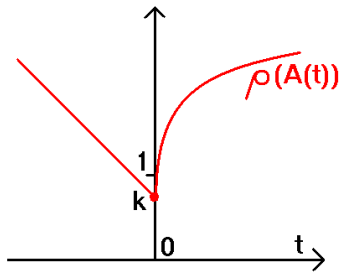
where  $\Lambda(A) = \{\text{eigenvalues of } A\}$  is the **spectrum**.

**Example**

The spectral radius of

$$A(t) = \begin{bmatrix} k & 1 \\ t & k-t \end{bmatrix}$$

(with  $k$  slightly less than 1)  
is minimized at  $t = 0$ .



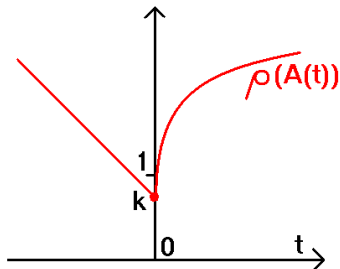
### 13. ROBUSTNESS AND TRANSIENT PEAKS

But

$$A(0) = \begin{bmatrix} k & 1 \\ 0 & k \end{bmatrix}$$

may be unsatisfactory.

**Difficulty I:**  $\rho(A(t))$  is highly sensitive to perturbation at  $t = 0$  (nonlipschitz).



**Difficulty II:** The trajectory  $\{A(0)^n\}$  has a big **transient peak:**

$$\begin{bmatrix} \frac{n}{n+1} & 1 \\ 0 & \frac{n}{n+1} \end{bmatrix}^n \sim e^{-1} \begin{bmatrix} 1 & n+1 \\ 0 & 1 \end{bmatrix} \quad \text{for large } n.$$

One difficulty is the **multiple eigenvalue**. But this is **typical** at optimal solutions of spectral radius minimization problems. (Burke/Lewis/Overton 2001)

## 14. PSEUDOSPECTRA

(Following Trefethen...)

A powerful tool to visualize robust properties of eigenvalues:

$$\Lambda_\epsilon(A) = \bigcup_{\|X-A\| \leq \epsilon} \Lambda(X) = \{z \in \mathbf{C} : \sigma_{\min}(A - zI) \leq \epsilon\},$$

where  $\sigma_{\min}$  is the **smallest singular value**.

Pseudospectra resolve Difficulty II (transient peaks)...

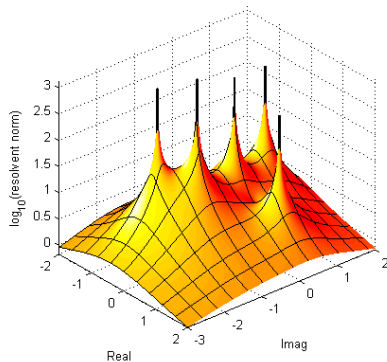
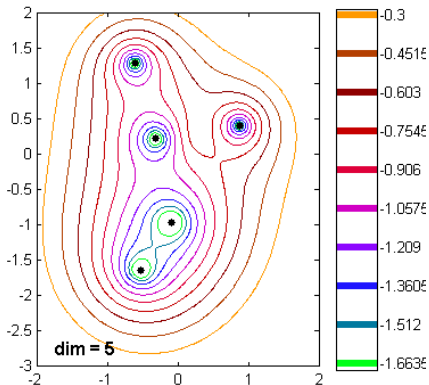
### **Kreiss Matrix Theorem (1962)**

$A^n < K\rho^n$  for all  $n$ , with  $K$  not too large  $\Leftrightarrow$   
 $\max\{|\lambda| : \lambda \in \Lambda_\epsilon(A)\} < \rho$ , with  $\epsilon$  not too small.

Analogously, in **continuous time**,  $e^{At} \rightarrow 0$  with peaks not too large when  $\Lambda_\epsilon(A)$  lies in the left halfplane for  $\epsilon$  not too small.

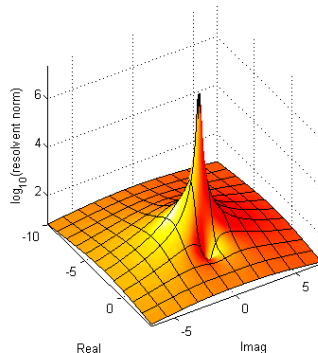
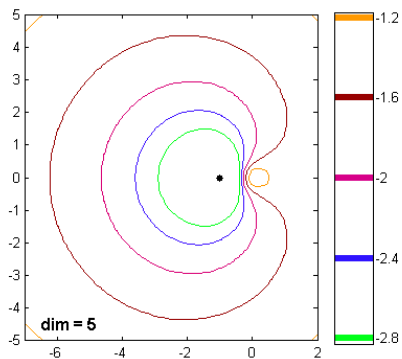
## 15. EXAMPLES

Pseudospectra for a random  $5 \times 5$  triangular complex matrix, plotted by **T. Wright's EigTool**:

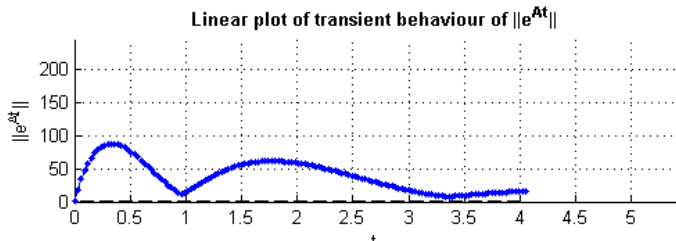


**Demmel's example:**  $A = - \begin{bmatrix} 1 & 5 & 5^2 & 5^3 & 5^4 \\ 0 & 1 & 5 & 5^2 & 5^3 \\ 0 & 0 & 1 & 5 & 5^2 \\ 0 & 0 & 0 & 1 & 5 \\ 0 & 0 & 0 & 0 & 1 \end{bmatrix} \dots$

## 16. A NEARLY UNSTABLE MATRIX



Notice  $\Lambda_{.01}(A)$  extends outside the left halfplane:  
some “unstable”  $X$  satisfies  $\|X - A\| \leq .01$ .



## 17. COMPUTING WITH PSEUDOSPECTRA

We can also **compute** pseudospectral quantities like

$$\alpha_\epsilon(A) = \max\{\operatorname{Re} \lambda : \lambda \in \Lambda_\epsilon(A)\}.$$

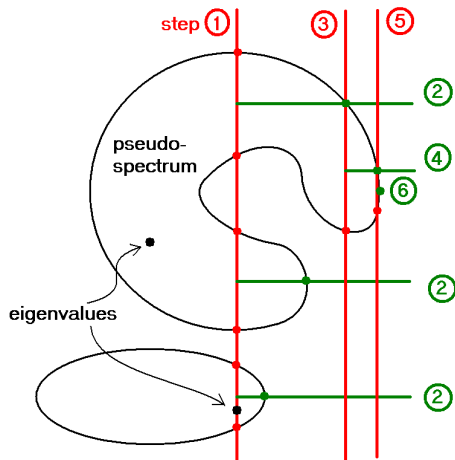
Key pseudospectral properties:

- Components all contain eigenvalues (by maximum modulus principle)
- Intersections with lines (or circles) easy to compute by a Hamiltonian eigensolver.

Hence a fast, accurate, robust **criss-cross algorithm** for  $\alpha_\epsilon$ :

globally and quadratically convergent, available in `eigtool`.

Returns  $\nabla \alpha_\epsilon$  (when it exists)  $\rightarrow$  nonsmooth **gradient sampling** for **optimizing**  $\alpha_\epsilon$  (Burke/Lewis/Overton 2003).



## 18. LIPSCHITZ BEHAVIOR

**Difficulty I?**  $A \mapsto \Lambda(A)$  isn't locally **Lipschitz**: no  $k$  satisfies

$$d(\Lambda(X), \Lambda(Y)) \leq k \|X - Y\| \quad \text{for all } X, Y \text{ near } A,$$

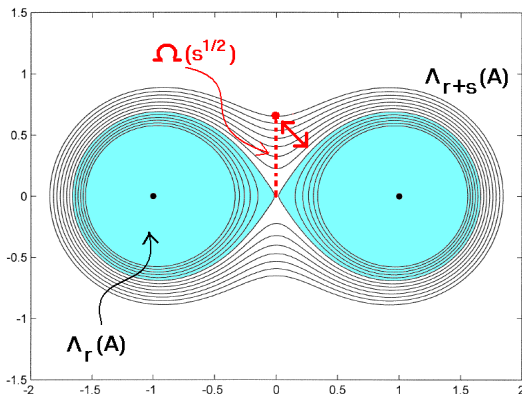
where  $d$  is Hausdorff distance between  $U, V \subset \mathbf{C}$ .

The pseudospectral map  $A \mapsto \Lambda_\epsilon(A)$  can also be nonlipschitz:

$$A = \begin{bmatrix} 1 & 1 \\ 0 & -1 \end{bmatrix}$$

$$r = \frac{\sqrt{5} - 1}{2}.$$

**But...**



**Theorem** Typically (Arnold 1971)  $A$  has all eigenspaces one-dimensional. Then,  $\Lambda_\epsilon$  is Lipschitz around  $A$  for all small  $\epsilon > 0$ .

## 19. CONTROLLABILITY

A **system** with **state**  $x \in \mathbf{C}^m$  and **control**  $u$ ,

$$\frac{dx}{dt} = Ax + Bu,$$

**controllable** if any endpoints can be interpolated by a control.

**Lemma (Hautus 1969)** A matrix-pair  $(A, B)$  is controllable

$$\Leftrightarrow \delta = \min\{\sigma_{\min}[A - zI, B] : z \in \mathbf{C}\} > 0.$$

Eising (1984) showed  $\delta$  is the **distance to uncontrollability**:

$$\delta = \min\{\|(X, Y)\| : (A + X, B + Y) \text{ is uncontrollable}\}$$

Polynomially computable (Gu 2000, 2005) but tough...

**Question (Trefethen)** How many components has the **rectangular pseudospectrum**

$$\{z \in \mathbf{C} : \sigma_{\min}[A - zI, B] \leq \epsilon\}?$$

## 20. CONNECTED COMPONENTS

**Problem:** For  $m$ -by- $n$   $P, Q$ , bound number of components

$$c = \#\{z : \sigma_{\min}(P + zQ) \leq \epsilon\}.$$

**Conjecture:**  $c \leq m$ . (Easy if  $m = 1$  or  $m = n$ .)

**Theorem**  $c \leq 2m^2 - m + 1$ . (Burke/Lewis/Overton 2004)

**Proof** Von Neumann-Wigner (1929) showed

$\{m$ -by- $m$  Hermitians with multiple eigenvalues $\}$

has codimension 3. Assume  $(P, Q)$  “typical”, so

$$\lambda_{\min}((P + zQ)(P + zQ)^*) \text{ simple } \forall z \in \mathbf{C}.$$

Milnor (1964) showed degree- $d$  polynomials  $p$  on  $\mathbf{R}^2$  satisfy

$$\#\{(x, y) : p(x, y) = 0\} \leq d(2d - 1),$$

Result follows by applying this to  $p : \mathbf{C} \cong \mathbf{R}^2 \rightarrow \mathbf{R}$  given by

$$p(z) = \det((P + zQ)(P + zQ)^* - \epsilon^2 I).$$

General case follows by perturbation.

## 21. SUMMARY

- Hyperbolicity is elegant for primal convex optimization.
  - But no apparent duality theory;
  - Is it really more general than semidefinite programming?
- Semisimple Lie theory gives a broad duality framework:
  - Fenchel conjugates;
  - convex and nonconvex subdifferentials.
- Spectral radius minimization leads to
  - multiple eigenvalues,
  - nonrobust solutions,
  - transient peaks.
- Pseudospectral optimization circumvents these (by Kreiss).
- Distance to uncontrollability is tractable.