

Lecture 15:Today:

- Stochastic DP.

- So far, we have been looking at deterministic Dynamic programming problems. This means that given the current state, and the decision, the state at the next stage is completely determined.

Ex: If the initial inventory at the beginning of period 1 is 0, and we produce $x_1 = 5$ units, then the initial inventory at the beginning of period 2 is completely determined: It is $I + x_1 - d_1$, where $d_1 =$ demand in period 1, known.

- In Stochastic dynamic programming problems, however, given the current state and a decision, the state at the next stage is not completely determined. Instead, there is a probability distribution that specifies a set of possible states at the next stage.

EX A betting game.

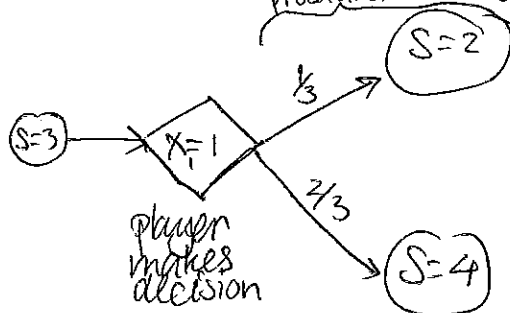
- Player starts with three chips.
- The game has three plays
- At each play, the player can bet any number of her chips
 - with probability $\frac{1}{3}$, the player loses # chips that are bet
 - with probability $\frac{2}{3}$, the player wins # chips that are bet.
- Goal: to finish the game with (at least) 5 chips.

• First, let's do a small example to see how this game is played.

Ex: Start with $s=3$ chips. \rightarrow can bet 0, 1, 2, or 3 chips.

Play #1: bet $x_1 = 1$ chip.

Then, will end up with 2 chips with prob = $\frac{1}{3}$, and
will end up with 4 chips with prob = $\frac{2}{3}$.
probabilistic outcome.



Note: Must wait to see the outcome before moving on to
... Play #2.

Suppose: $S=2$. \rightarrow can bet 0, 1, or 2 chips.

Play #2: bet $x_2 = 1$ chip.

Then will end up w/ 1 chip w/ prob = $\frac{1}{3}$
3 chips w/ prob = $\frac{2}{3}$.

Then, suppose $S=3$

Play #3: bet $x_3 = 2$ chips.

Then will end up with 1 chip w/ prob = $\frac{1}{3}$
3 chips w/ prob = $\frac{2}{3}$.

\rightarrow Can "think of this as a 2-player game: player vs. "nature"
where "nature" is a probabilistic player.

• A Dynamic programming approach :

• Stage $k \leftrightarrow$ play # k .

• State \leftrightarrow remaining # coins at beginning of the stage.
 $\rightarrow S_1 = \{3\}$, $S_2 = \{0, 1, \dots\}$, $S_3 = \{0, \dots, 5\}$.

• At stage k , state s ,

decide: $X_k =$ # coins to bet at stage k .

$X_k \in \{0, 1, \dots, \min\{s, 5-s\}\}$.

Note: X_k can be any # from $0, 1, \dots, s$. $\left. \begin{array}{l} \text{s.t. } S - X_k \geq 0 \text{ and } S + X_k \leq 5. \end{array} \right\} \begin{array}{l} \text{so } X_k \leq s \text{ and} \\ X_k \leq 5 - s. \end{array}$

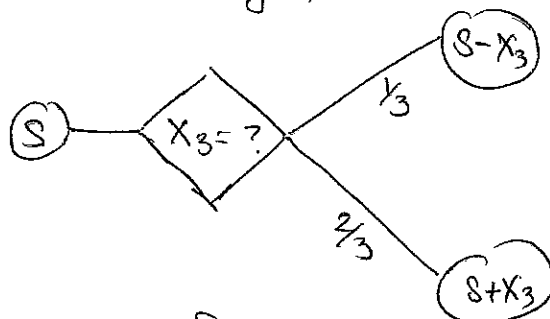
• Optimization function:

$f_k^*(s) =$ max probability of having at least 5 coins at the end, given that s coins remain at the beginning of stage k .



Note: At stage k , we consider plays $k, k+1, \dots$, end.

• Suppose $k=3 =$ last stage, with s coins remaining:



- Make decision $X_3 \in \{0, 1, \dots, \min\{s, 5-s\}\}$.

- Since this is the last stage, then the probability of winning the game is simple to compute:

• Check if $s - X_3 \geq 5$ or $s + X_3 \geq 5$.

• Then, prob of winning the game is
$$u = \begin{cases} 1 & \text{if } s - X_3 \geq 5, s + X_3 \geq 5 \\ 2/3 & \text{if } s - X_3 < 5, s + X_3 \geq 5 \\ 0 & \text{if } s - X_3 < 5, s + X_3 < 5 \end{cases}$$

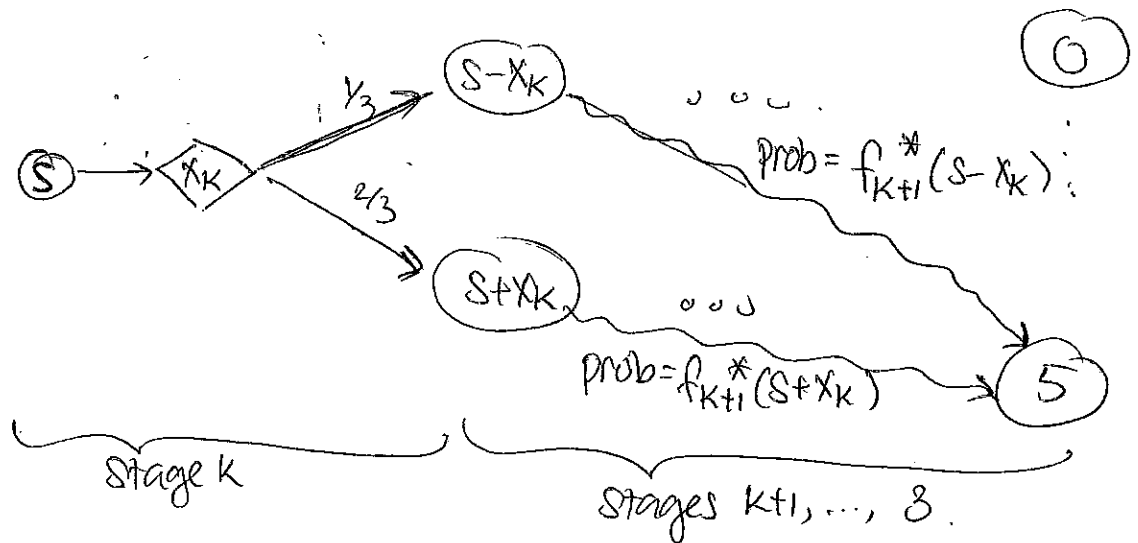
→ Choose x_3 that maximizes the probability of winning.

$f_3^*(0) = 0$	$x_3^* = 0$
$f_3^*(1) = 0$	$x_3^* = 0, 1$
$f_3^*(2) = 0$	$x_3^* = 0, 1, 2$
$f_3^*(3) = 2/3$	$x_3^* = 2$
$f_3^*(4) = 2/3$	$x_3^* = 1$
$f_3^*(5) = 1$	$x_3^* = 0$

These are the boundary conditions.

• Recurrence relation:

$$f_k^*(s) = \max_{\substack{x_k=0,1,\dots, \\ \min\{s, 5-s\}}} \left\{ \begin{array}{l} \text{prob that we'll have 5 coins at end of 3rd stage} \\ \text{if we start at beginning of stage } k, \text{ and we bet } x_k \text{ coins.} \end{array} \right.$$



So, Starting w/ s coins at stage k , if we bet x_k coins, the prob. of having 5 coins at the end is:

$$\frac{1}{3} \cdot f_{k+1}^*(s - x_k) + \frac{2}{3} f_{k+1}^*(s + x_k). \quad (1)$$

So, choose x_k that maximizes (1):

$$f_k^*(s) = \max_{\substack{x_k \\ \in \{0, \dots, \min\{s, 5-s\}\}}} \left\{ \frac{1}{3} f_{k+1}^*(s - x_k) + \frac{2}{3} f_{k+1}^*(s + x_k) \right\}.$$

• Computation:

Stage 3 \rightarrow boundary conditions.

Stage 2:

$$f_2^*(0) = \frac{1}{3}f_3^*(0) + \frac{2}{3}f_3^*(0) = 0 \quad x_2^* = 0.$$

$$\begin{aligned} f_2^*(1) &= \max_{x_2 \in \{0,1\}} \left\{ \frac{1}{3}f_3^*(1-x_2) + \frac{2}{3}f_3^*(1+x_2) \right\} \\ &= \max \left\{ \underbrace{\frac{1}{3}f_3^*(1) + \frac{2}{3}f_3^*(1)}_{x_2=0}, \underbrace{\frac{1}{3}f_3^*(0) + \frac{2}{3}f_3^*(2)}_{x_2=1} \right\} \\ &= 0 \quad x_2^* = 0, 1. \end{aligned}$$

$$\begin{aligned} f_2^*(2) &= \max \left\{ \frac{1}{3}f_3^*(2) + \frac{2}{3}f_3^*(2), \frac{1}{3}f_3^*(1) + \frac{2}{3}f_3^*(3), \right. \\ &\quad \left. \frac{1}{3}f_3^*(0) + \frac{2}{3}f_3^*(4) \right\} \\ &= 4/9 \quad x_2^* = 1 \text{ or } 2. \end{aligned}$$

$$\begin{aligned} f_2^*(3) &= \max \left\{ \frac{1}{3}f_3^*(3) + \frac{2}{3}f_3^*(3), \frac{1}{3}f_3^*(2) + \frac{2}{3}f_3^*(4), \right. \\ &\quad \left. \frac{1}{3}f_3^*(1) + \frac{2}{3}f_3^*(5) \right\} = \frac{2}{3} \quad x_2^* = 2. \end{aligned}$$

$$\begin{aligned} f_2^*(4) &= \max \left\{ \frac{1}{3}f_3^*(4) + \frac{2}{3}f_3^*(4), \frac{1}{3}f_3^*(3) + \frac{2}{3}f_3^*(5) \right\} \\ &= \frac{1}{3} \cdot \frac{2}{3} + \frac{2}{3} = \frac{8}{9} \quad x_2^* = 1. \end{aligned}$$

$$f_2^*(5) = \frac{1}{3}f_3^*(5) + \frac{2}{3}f_3^*(5) = 1 \quad x_2^* = 0.$$

Stage 1:

$$f_1^*(0) = 0 \quad X_1^* = 0$$

$$f_1^*(1) = \max \{ f_2^*(1), \frac{1}{3}f_2^*(0) + \frac{2}{3}f_2^*(1) \} = 0 \quad X_1^* = 0$$

$$f_1^*(2) = \max \{ f_2^*(2), \frac{1}{3}f_2^*(1) + \frac{2}{3}f_2^*(3), \frac{1}{3}f_2^*(0) + \frac{2}{3}f_2^*(4) \}$$

$$= \frac{16}{27} \quad X_1^* = 2$$

$$f_1^*(3) = \max \{ f_2^*(3), \frac{1}{3}f_2^*(2) + \frac{2}{3}f_2^*(4), \frac{1}{3}f_2^*(1) + \frac{2}{3}f_2^*(5) \}$$

$$= \max \{ \frac{2}{3}, \frac{1}{3} \cdot \frac{4}{9} + \frac{2}{3} \cdot \frac{8}{9}, \frac{2}{3} \}$$

$$= \frac{20}{27}, \quad X_1^* = 1$$

$$f_1^*(4) = \max \{ f_2^*(4), \frac{1}{3}f_2^*(3) + \frac{2}{3}f_2^*(5) \}$$

$$= \max \{ \frac{8}{9}, \frac{1}{3} \cdot \frac{2}{3} + \frac{2}{3} \cdot 1 \} = \frac{8}{9} \quad X_1^* = 0, 1$$

$$f_1^*(5) = 1 \quad X_1^* = 0$$

DP Table: $f_k^*(s)$

	<u>Stage 1</u>		<u>Stage 2</u>		<u>Stage 3</u>	
$s = 0$	0	$X_1^* = 0$	0	$X_2^* = 0$	0	$X_3^* = 0$
1	0	$X_1^* = 0$	0	$X_2^* = 0, 1$	0	$X_3^* = 0, 1$
2	$\frac{16}{27}$	$X_1^* = 2$	$\frac{4}{9}$	$X_2^* = 1, 2$	0	$X_3^* = 0, 1, 2$
3	$\frac{20}{27}$	$X_1^* = 1$	$\frac{2}{3}$	$X_2^* = 2$	$\frac{2}{3}$	$X_3^* = 2$
4	$\frac{8}{9}$	$X_1^* = 0, 1$	$\frac{8}{9}$	$X_2^* = 1$	$\frac{2}{3}$	$X_3^* = 1$
5	1	$X_1^* = 0$	1	$X_2^* = 0$	1	$X_3^* = 0$

In Game Theory "Lingo": this table gives you a "complete strategy" for maximizing your probability of winning the game. That is, at each play, you only need to see how many coins remain and play X_k^* according to the above table.