

# Lecture 11

# Lecture 11

*Previously in Opt 2 ...*

# Example: Inventory planning problem

Period (k)	$d_k$	$c_k$	$F_k$	$h_k$
1	10	3	5	0.2
2	40	2	20	0.3
3	20	4	10	0.5
4	50	3	10	0.8
5	0			

# Example: Inventory planning problem

## DP Formulation

### 1. Specify stages

stage  $k \leftrightarrow$  period  $k$

### 2. Specify states at each stage $k$

State  $I \leftrightarrow$  inventory level  $I$

(that is,  $I$  is the quantity in the inventory)

### 3. Specify allowable decisions at each state $I$ and stage $k$

values  $x_k$  such that demand  $d_k$  can be satisfied

i.e.,  $x_k$  such that  $I + x_k \geq d_k$

# Example: Inventory planning problem

## DP Formulation

### 1. Specify stages

Stages:  $k = 1, 2, 3, 4, 5$

### 2. Specify states at each stage $k$

$$S_k = \{0, 10, 20, \dots, 120\}$$

( $S_k$  is the set of possible inventory levels at stage  $k$ )

### 3. Specify allowable decisions at each state $I$ and stage $k$

$$Q_{I,k} = \{x \text{ in } \{0, 10, \dots, 120\} \mid I + x \geq d_k\}$$

# Example: Inventory planning problem

## DP Formulation

4. Word-description of optimal function to be solved at state  $I$  in stage  $k$

$f_k^*(I)$  = the minimum cost for satisfying demands in periods  $k, k+1, \dots, T, T+1$

5. Boundary conditions

$$f_{T+1}^*(I) = 0 \text{ for all } I$$

6. Recurrence relation

$$f_k^*(I) = \min_{x_k \text{ in } Q_{k,I}} \{c_k x_k + f_k 1_{\{x_k > 0\}} + h_k (I + x_k - d_k) + f_{k+1}^*(I + x_k - d_k)\}$$

# Example: Inventory planning problem

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$f_k^*(I)$  = the minimum cost for satisfying demands in periods  $k, k+1, \dots, T, T+1$

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$$f_{T+1}^*(I) = 0 \text{ for all } I$$

6. Recurrence relation

$$\underbrace{f_k^*(I)}_{\text{Min cost in periods } k \text{ to } T+1 \text{ given that inventory at start is } I} = \min_{x_k \text{ in } Q_{k,I}} \{c_k x_k + f_k 1_{\{x_k > 0\}} + h_k (I + x_k - d_k) + f_{k+1}^*(I + x_k - d_k)\}$$

Min cost in periods  $k$  to  $T+1$  given that inventory at start is  $I$

# Example: Inventory planning problem

## DP Formulation

4. Word-description of optimal function to be solved at state  $I$  in stage  $k$

$f_k^*(I)$  = the minimum cost for satisfying demands in periods  $k, k+1, \dots, T, T+1$

5. Boundary conditions

$$f_{T+1}^*(I) = 0 \text{ for all } I$$

6. Recurrence relation

$$f_k^*(I) = \min_{x_k \text{ in } Q_{k,I}} \left\{ \underbrace{c_k x_k + f_k \mathbf{1}_{\{x_k > 0\}} + h_k (I + x_k - d_k)}_{\text{Cost incurred in period } k} + f_{k+1}^*(I + x_k - d_k) \right\}$$



# Example: Inventory planning problem

## DP Formulation

4. Word-description of optimal function to be solved at state  $I$  in stage  $k$

$f_k^*(I)$  = the minimum cost for satisfying demands in periods  $k, k+1, \dots, T, T+1$

5. Boundary conditions

$$f_{T+1}^*(I) = 0 \text{ for all } I$$

6. Recurrence relation

$$f_k^*(I) = \min_{x_k \text{ in } Q_{k,I}} \{ c_k x_k + f_k 1_{\{x_k > 0\}} + h_k (I + x_k - d_k) + \underbrace{f_{k+1}^*(I + x_k - d_k)}_{\text{Min cost in periods } k+1 \text{ to } T+1} \}$$

Min cost in periods  $k+1$  to  $T+1$

# Example: Inventory planning problem

## DP Formulation

4. Word-description of optimal function to be solved at state  $I$  in stage  $k$

$f_k^*(I)$  = the minimum cost for satisfying demands in periods  $k, k+1, \dots, T, T+1$

5. Boundary conditions

$$f_{T+1}^*(I) = 0 \text{ for all } I$$

6. Recurrence relation

$$f_k^*(I) = \min_{\underbrace{x_k \text{ in } Q_{k,I}}_{\text{Old inventory level}}} \{ c_k x_k + f_k 1_{\{x_k > 0\}} + h_k (I + x_k - d_k) + f_{k+1}^*(\underbrace{I + x_k - d_k}_{\text{New inventory level}}) \}$$

Old inventory level


New inventory level

More generally, recursive relations look like:

$$f_k^*(\textit{current state}) = \min_{x_k \text{ in } Q_{k,I}} \left\{ \text{value/cost due to decision } x_k + f_{k+1}^*(\textit{new state due to decision } x_k) \right\}$$

More generally, recursive relations look like:

(Or another operation)

$$f_k^*(\textit{current state}) = \min_{x_k \text{ in } Q_{k,I}} \left\{ \text{value/cost due to decision } x_k + f_{k+1}^*(\textit{new state due to decision } x_k) \right\}$$


# Example: Inventory planning problem

## DP Formulation

### 7. Computation

Summary of DP Computation:

Possible States (i)	Stage 5	Stage 4		Stage 3		Stage 2		Stage 1	
	$f^*_5(i)$	$f^*_4(i)$	$x^*_4$	$f^*_3(i)$	$x^*_3$	$f^*_2(i)$	$x^*_2$	$f^*_1(i)$	$x^*_1$
0	0					286	110		
10	0					266	100		
20	0					246	90		
30	0					226	80		
40	0								
50	0								
60	0								
70	0								
80	0								
90	0					100	0		
100	0					78	0		
110	0					46	0		

# Example: Inventory planning problem

## DP Formulation

### 8. Trace back to find optimal solution

Summary of DP Computation:

Possible States (i)	Stage 5	Stage 4		Stage 3		Stage 2		Stage 1	
	$f^*_5(i)$	$f^*_4(i)$	$x^*_4$	$f^*_3(i)$	$x^*_3$	$f^*_2(i)$	$x^*_2$	$f^*_1(i)$	$x^*_1$
0	0					286	110		
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100	0					78	0		
110	0					46	0		

# Example: Inventory planning problem

## DP Formulation

### 8. Trace back to find optimal solution

Possible States (i)	Stage 5	Stage 4		Stage 3		Stage 2		Stage 1	
	$f^*_5(i)$	$f^*_4(i)$	$x^*_4$	$f^*_3(i)$	$x^*_3$	$f^*_2(i)$	$x^*_2$	$f^*_1(i)$	$x^*_1$
0	0	160	50	250	20	286	110	321	10
10	0	130	40	210	10	266	100		
20	0	100	30	160	0	246	90		
30	0	70	20	135	0	226	80		
40	0	40	10	110	0	206	70		
50	0	0	0	85	0	186	60		
60	0			60	0	166	0 or 50		
70	0			25	0	144	0		
80	0					122	0		
90	0					100	0		
100	0					78	0		
110	0					46	0		

i>clicker question



Q: What is the optimal quantity to produce at each time period?

Possible States (i)	Stage 5	Stage 4		Stage 3		Stage 2		Stage 1	
	$f^*_5(i)$	$f^*_4(i)$	$x^*_4$	$f^*_3(i)$	$x^*_3$	$f^*_2(i)$	$x^*_2$	$f^*_1(i)$	$x^*_1$
0	0	160	50	250	20	286	110	321	10
10	0	130	40	210	10	266	100		
20	0	100	30	160	0	246	90		
30	0	70	20	135	0	226	80		
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60	0			60	0	166	0 or 50		
70	0			25	0	144	0		
80	0					122	0		
90	0					100	0		
100	0					78	0		
110	0					46	0		

- A.  $x_1 = 10, x_2 = 110, x_3 = 20, x_4 = 50$
- B.  $x_1 = 10, x_2 = 0, x_3 = 0, x_4 = 0$
- C.  $x_1 = 10, x_2 = 110, x_3 = 0, x_4 = 0$
- D.  $x_1 = 10, x_2 = 40, x_3 = 20, x_4 = 50$
- E.  $x_1 = 10, x_2 = 0, x_3 = 110, x_4 = 0$



Q: What is the optimal quantity to produce at each time period?

Summary of DP Computation:

Possible States (i)	Stage 5	Stage 4		Stage 3		Stage 2		Stage 1	
	$f^*_5(i)$	$f^*_4(i)$	$x^*_4$	$f^*_3(i)$	$x^*_3$	$f^*_2(i)$	$x^*_2$	$f^*_1(i)$	$x^*_1$
0	0	160	50	250	20	286	110	321	10
10	0	130	40	210	10	266	100		
20	0	100	30	160	0	246	90		
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70	0			25	0	144	0		
80	0					122	0		
90	0					100	0		
100	0					78	0		
110	0					46	0		

- A.  $x_1 = 10, x_2 = 110, x_3 = 20, x_4 = 50$
- B.  $x_1 = 10, x_2 = 0, x_3 = 0, x_4 = 0$
- C.  $x_1 = 10, x_2 = 110, x_3 = 0, x_4 = 0$
- D.  $x_1 = 10, x_2 = 40, x_3 = 20, x_4 = 50$
- E.  $x_1 = 10, x_2 = 0, x_3 = 110, x_4 = 0$

# Example: Inventory planning problem

## DP Formulation

### 8. Trace back to find optimal solution

Summary of DP Computation:

Possible States (i)	Stage 5	Stage 4		Stage 3		Stage 2		Stage 1	
	$f^*_5(i)$	$f^*_4(i)$	$x^*_4$	$f^*_3(i)$	$x^*_3$	$f^*_2(i)$	$x^*_2$	$f^*_1(i)$	$x^*_1$
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# DP Computation using AMPL

As usual, need:

1. Model file
2. Data file
3. Script file
  - Optional, but this can save you time from having to re-enter the following commands many times while debugging/re-running:

```
model myDP.mod;  
data myDP.dat;  
solve;  
(etc.)
```

# DP Computation using AMPL

## 1. Model file

- We are not going to model a linear program here, so we won't have:

```
var  
maximize ...  
subject to ...
```

- Will use only sets and parameters:
  - The values of some sets/parameters are supplied by the data file
  - The values of some other sets/parameters are **computed within** the model file; not supplied by the data file

# DP Computation using AMPL

## 2. Data file

- Same as before!
- Here, you specify values of parameters that are declared in the model file.

# DP Computation using AMPL

## 2. Data file

- Same as before!
- Here, you specify values of parameters that are declared in the model file.

## 3. Script file

- Same as before!
- If you haven't been using scripts, you can start now!



# DP Computation using AMPL

1. inventory.mod
2. inventory.dat
3. inventoryScript.txt

# inventory.mod

```
param T;                                # number of periods
param c {k in 1..T};                    # cost per unit in period k
param F {k in 1..T};                    # fixed cost in period k
param h {k in 1..T};                    # holding cost for inventory from k-1 to k
param d {k in 1..T};                    # demand in period k

param MaxQuantity := sum{k in 1..T} d[k];

set allowableDecisions {k in 1..T, I in 0..MaxQuantity} :=
  {x in 0..MaxQuantity: d[k] <= I + x <= MaxQuantity};
  # quantity to produce at stage k, state I (inventory level)

param f {k in 1..T+1, I in 0..MaxQuantity} := # min cost in periods k through T+1
  if k=T+1 then 0 # no more demand to consider
  else
    min{x in allowableDecisions[k, I]}
    (c[k]*x + F[k]*(if x <> 0 then 1 else 0) + h[k]*(I + x - d[k]) + f[k+1, I + x - d[k]]);

set opt {k in 1..T, I in 0..MaxQuantity} := # optimal decisions
  {x in allowableDecisions[k, I]: f[k, I] = c[k]*x + F[k]*(if x <> 0 then 1 else 0) +
    h[k]*(I + x - d[k]) + f[k+1, I + x - d[k]]};

param DPvalue := f[1, 0]; # compute the optimal value
```

# inventory.mod

```
param T;                # number of periods
param c {k in 1..T};    # cost per unit in period k
param F {k in 1..T};    # fixed cost in period k
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  else
    min{x in allowableDecisions[k, I]} (
      c[k]*x + F[k]*(if x <> 0 then 1 else 0) + h[k]*(I + x - d[k]) + f[k+1, I + x - d[k]]);

set opt {k in 1..T, I in 0..MaxQuantity} := # optimal decisions
  {x in allowableDecisions[k, I]: f[k, I] = c[k]*x + F[k]*(if x <> 0 then 1 else 0) +
    h[k]*(I + x - d[k]) + f[k+1, I + x - d[k]]};

param DPvalue := f[1, 0]; # compute the optimal value
```

Values are provided by inventory.dat

# inventory.mod

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param F {k in 1..T};                    # fixed cost in period k
param h {k in 1..T};                    # holding cost for inventory from k-1 to k
param d {k in 1..T};                    # demand in period k

param MaxQuantity := sum{k in 1..T} d[k]; } Max quantity considered

set allowableDecisions {k in 1..T, I in 0..MaxQuantity} :=
  {x in 0..MaxQuantity: d[k] <= I + x <= MaxQuantity};
  # quantity to produce at stage k, state I (inventory level)

param f {k in 1..T+1, I in 0..MaxQuantity} := # min cost in periods k through T+1
  if k=T+1 then 0 # no more demand to consider
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param DPvalue := f[1, 0]; # compute the optimal value
```

} Set of allowable decisions at each stage k, state I

# inventory.mod

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    h[k]*(I + x - d[k]) + f[k+1, I + x - d[k]]};

param DPvalue := f[1, 0]; # compute the optimal value
```

Computation of  $f_k^*(I)$   
at each stage  $k$ , each  
state  $I$

# inventory.mod

```
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  else
    min{x in allowableDecisions[k, I]}
    (c[k]*x + F[k]*(if x <> 0 then 1 else 0) + h[k]*(I + x - d[k]) + f[k+1, I + x - d[k]]);

set opt {k in 1..T, I in 0..MaxQuantity} := # optimal decisions
  {x in allowableDecisions[k, I]: f[k, I] = c[k]*x + F[k]*(if x <> 0 then 1 else 0) +
    h[k]*(I + x - d[k]) + f[k+1, I + x - d[k]]};

param DPvalue := f[1, 0]; # compute the optimal value
```

Keeping track of optimal decision,  $x_k^*$ , at each stage  $k$ , each state  $I$

# inventory.mod

```
param T;                                # number of periods
param c {k in 1..T};                    # cost per unit in period k
param F {k in 1..T};                    # fixed cost in period k
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set opt {k in 1..T, I in 0..MaxQuantity} := # optimal decisions
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    h[k]*(I + x - d[k]) + f[k+1, I + x - d[k]]};

param DPvalue := f[1, 0]; # compute the optimal value
```

Determining the optimal  
cost:  $f_1^*(0)$



# inventory.dat

```
param T := 4; # number of periods
```

```
param: d c F h :=
```

```
1      10    3    5 0.2
```

```
2      40    2  20 0.3
```

```
3      20    4  10 0.5
```

```
4      50    3  10 0.8;
```

# inventoryScript.txt

```
reset;  
model Opt2Models/Lec11/inventory.mod;  
data Opt2Models/Lec11/inventory.dat;  
display DPvalue;
```

# In the “sw” console:

```
sw: ampl
```

```
ampl: include Opt2Models/Lec11/inventoryScript.txt;
```

```
DPvalue = 321
```

```
ampl:
```